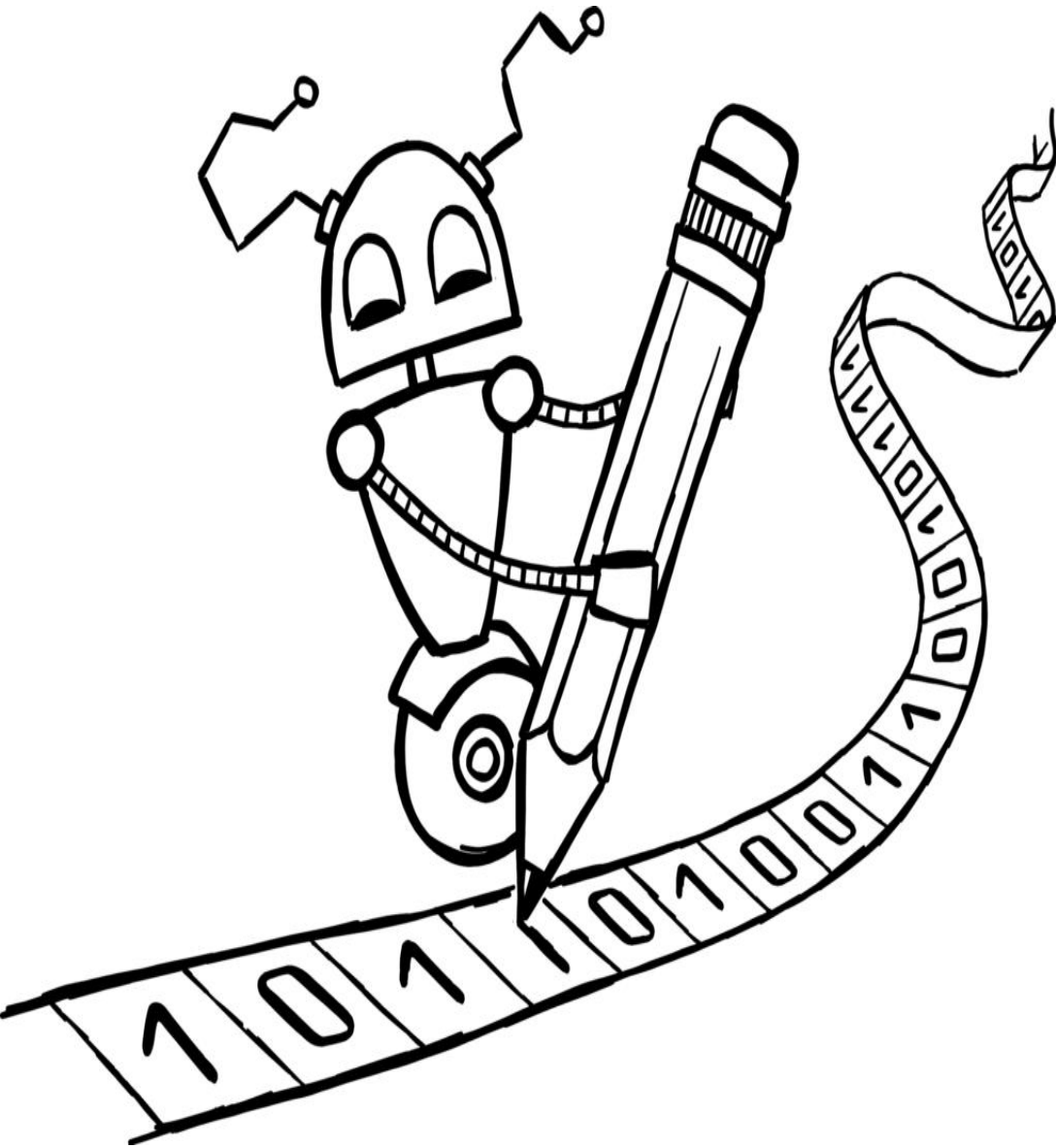




Computing

Analiza Algoritmulor

Doing nothing

	0	1	$\square$
q_1	$q_1, 0, -$	$q_1, 1, -$	$q_1, \square, -$

Writing 1 forever

	0	1	□
q_1	$q_1, 1, \rightarrow$	$q_1, 1, \rightarrow$	$q_1, 1, \rightarrow$

Flipping bits forever

	0	1	$\square$
q_1	$q_1, 1, \rightarrow$	$q_1, 0, \rightarrow$	$q_2, \square, \leftarrow$
q_2	$q_2, 1, \leftarrow$	$q_2, 0, \leftarrow$	$q_1, \square, \rightarrow$

Halting on some inputs

	0	1	$\square$
q_1	H, 0, $-$	$q_2, 1, \rightarrow$	$q_2, 1, \rightarrow$
q_2	$q_2, 1, \rightarrow$	$q_2, 1, \rightarrow$	$q_2, 1, \rightarrow$

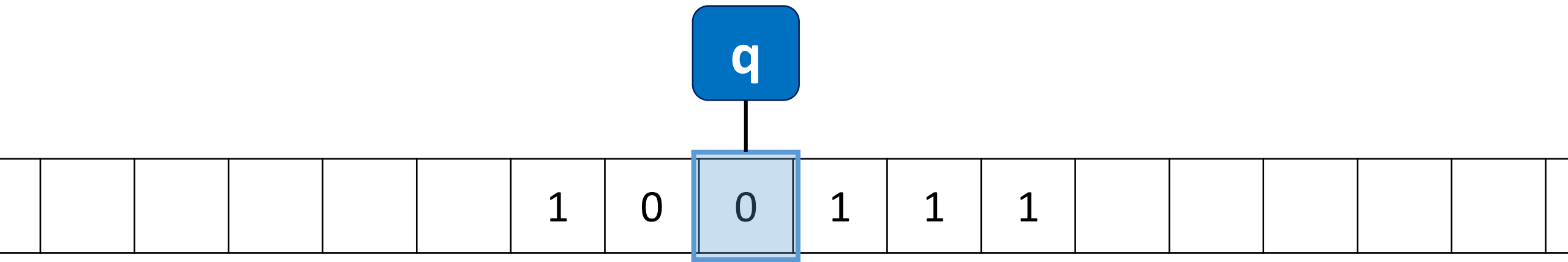
When a machine runs...

- ... it can ***halt*** in one of the final states: Y, N, H
- ... it can go on forever, ***without halting***

When a machine halts on input w ...

- ... in final state H , it *computes* a word v (the tape contents)
- ... in final state Y , it *accepts* w
- ... in final state N , it *rejects* w

Configurations



$(10, 0111, q)$

Computation history

$(\varepsilon, 110111, q)$

$(11011, 1, p)$

$(1, 10111, q)$

$(1101, 10, p)$

$(11, 0111, q)$

$(110, 100, p)$

$(110, 111, q)$

$(11, 0000, p)$

$(1101, 11, q)$

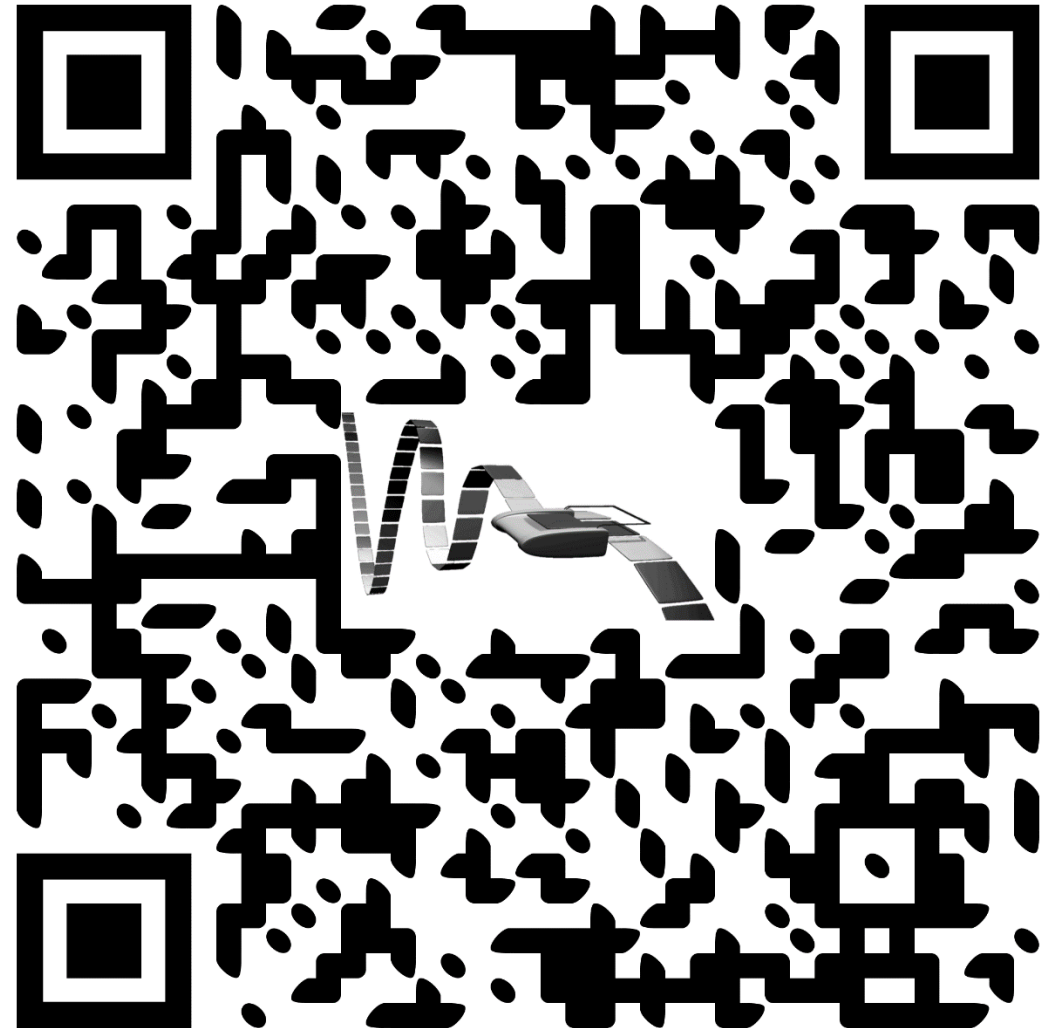
$(11, 1000, H)$

$(11011, 1, q)$

$(110111, \square, q)$

Moodle Quiz

Machine Configurations



Solving a problem

For a function problem f :

- If a machine M *always halts* with the *correct answer*, then " M **computes** f "

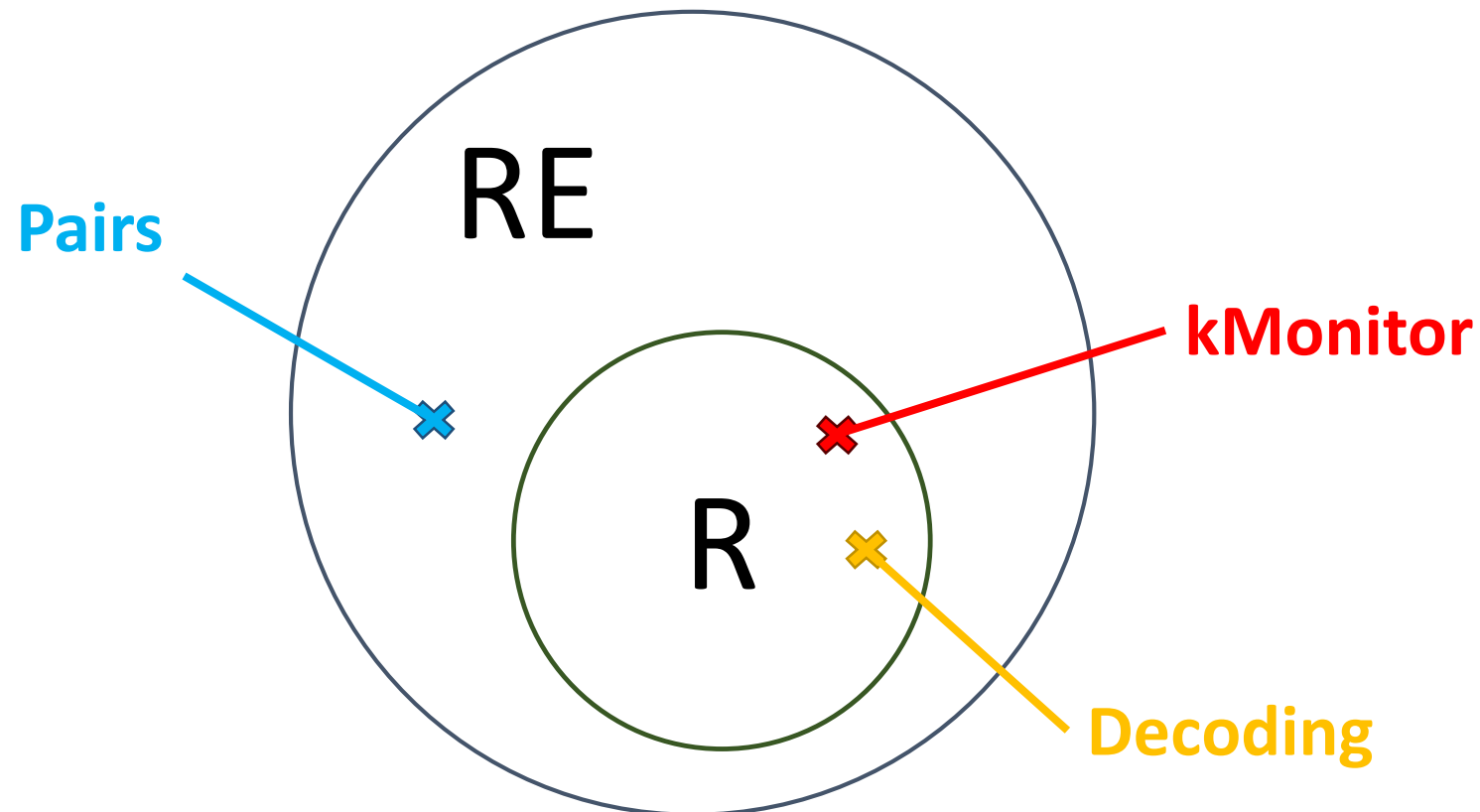
For a decision problem d :

- If a machine M *accepts* whenever d is TRUE, then " M **accepts** d "
- If a machine M *always halts* with the *correct answer*, then " M **decides** d "

The sets R, RE

$RE = \{f \mid f \text{ is acceptable}\}$

$R = \{f \mid f \text{ is decidable}\}$

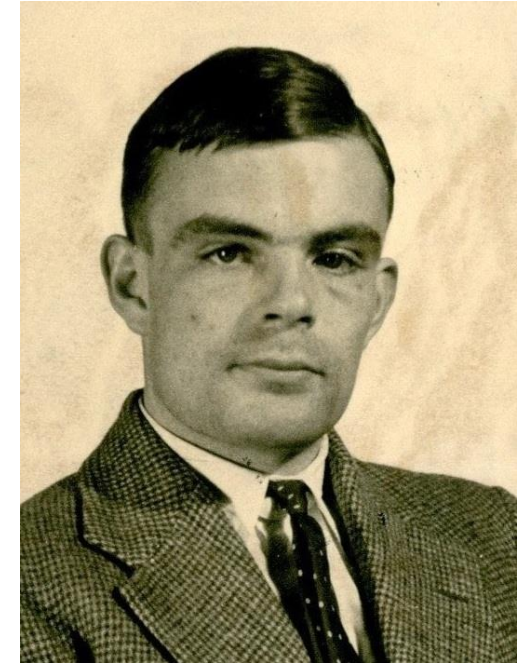


The Church-Turing Thesis



*Alonzo Church
(1903 – 1995)*

*“The problems solvable by algorithms
are exactly those
computable by Turing machines.”*



*Alan Turing
(1912 – 1954)*

All roads lead to R

Recursive functions

λ -calculus

Turing machines

Markov algorithms

RAM machines

Queue automata

Counter machines

Post canonical system

Wang B-machine

Uniform boolean circuits