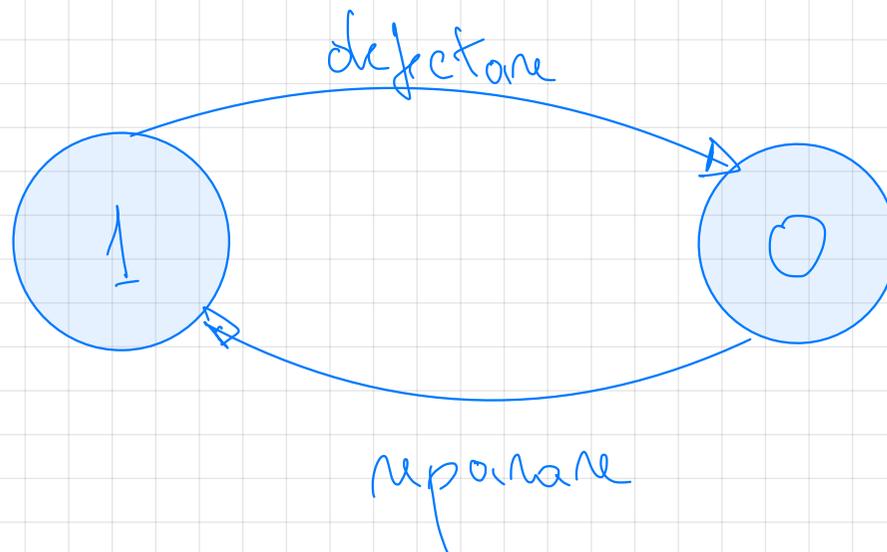
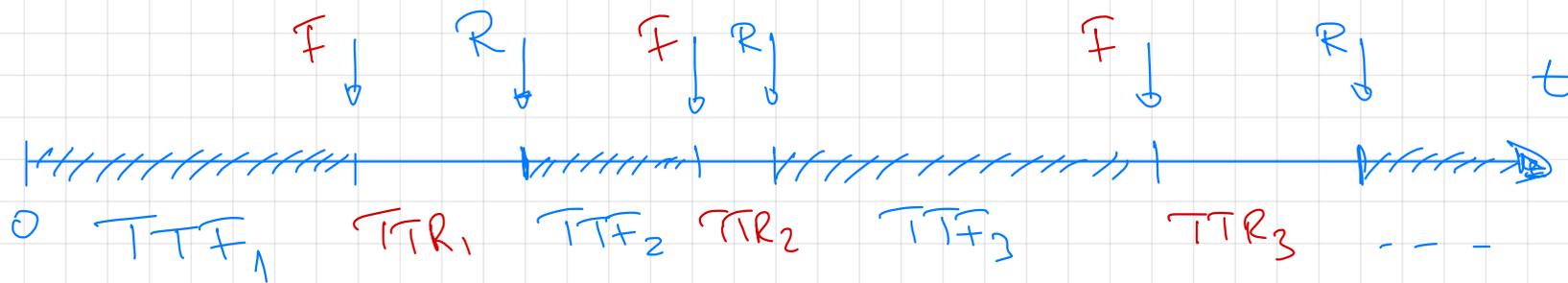


Fiabilitate și toleranță la defecte

$R(t)$ - fiabilitatea

$$R(t) = P(\tau > t \mid OK @ t=0)$$





$$MTBF = \sum_i \frac{TTF_i}{n}$$

$$MTR = \sum_i \frac{TR_i}{n}$$

Disponibilitate - $A(t)$

$$A = \frac{\sum_i TTF_i}{\sum_i TTF_i + \sum_i TR_i} = \frac{MTBF}{MTBF + MTR}$$

Availability (%)	Downtime / year	Downtime / month	Downtime / week
90% ("one nine")	36,5 days	72 h	16,8 h
99% ("2 nines")	3,65 days	7,2 h	1,68 h
99,9% ("3 nines")	8,76 h	43,2 min	10,1 min
99,99% ("4 nines")	52,56 min	4,32 min	1,01 min
99,999%	5,25 min	25,9 s	6,05 s
99,9999%	31,5 s	2,59 s	0,605 s

Probability theory 101

$$0 \leq P(A) \leq 1$$

$$P(\bar{A}) = 1 - P(A)$$

$P(A|B)$ - Prob A cond de B

$$P(A \cdot B) = P(A|B) \cdot P(B)$$

Doacă A și B sunt independente

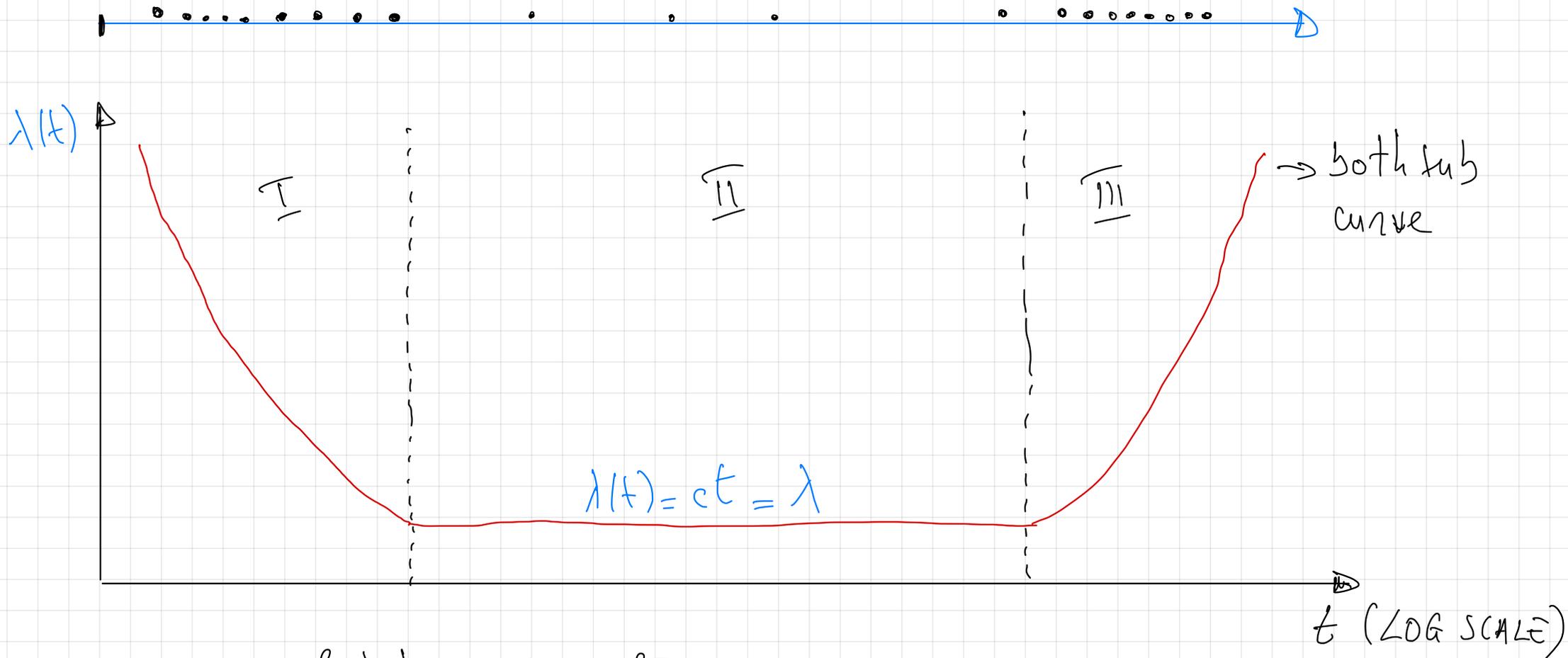
$$P(A \cdot B) = P(A) \cdot P(B)$$

$$P(A + B) = P(A) + P(B) - P(A \cdot B)$$

Doacă A și B sunt mutual exclusive $\Rightarrow P(A \cdot B) = P(B \cdot A) = 0$

$$P(A + B) = P(A) + P(B)$$

Failure Rate $\rightarrow \lambda(t)$ - intensitatea defectiunilor

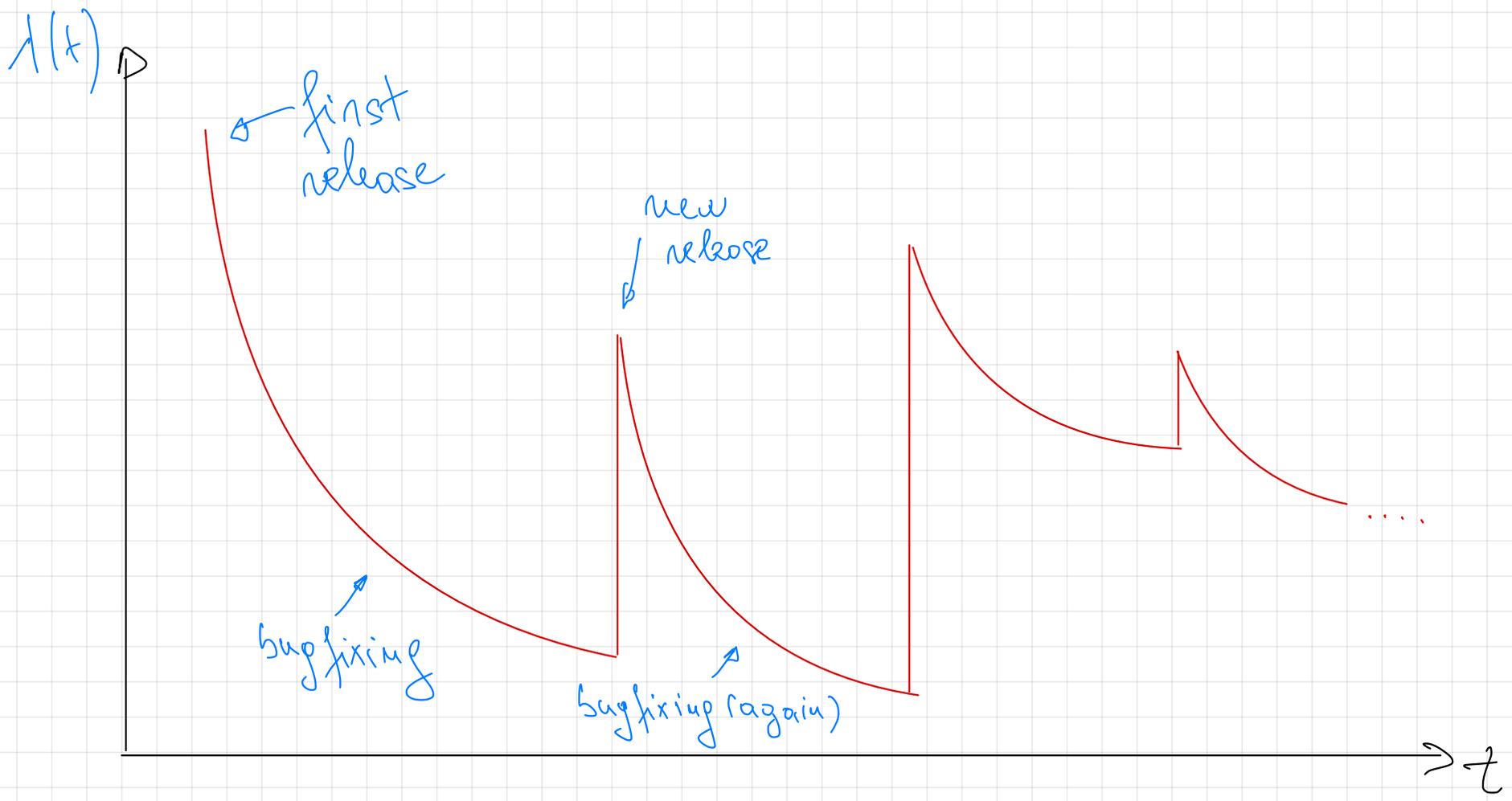


I - mortalitate infanțită

II - viață utilă

III - îmbătrânire

Failure rate - software



Câteva definiții noi:

$f(t)$ - funcție densitate probabilitate (pdf)

$F(t)$ - funcție cumulativă distribuție probabilitate

$$f(t) = \frac{dF(t)}{dt}$$

$$F(t) = \int_0^t f(z) dz$$

$$R(t) = 1 - F(t)$$

$$\lambda(t) = \frac{f(t)}{1 - F(t)}$$

$$\Rightarrow \lambda(t) = \frac{f(t)}{R(t)}$$

$$f(t) = \frac{dF(t)}{dt} = \frac{d(1 - R(t))}{dt} = - \frac{dR(t)}{dt}$$

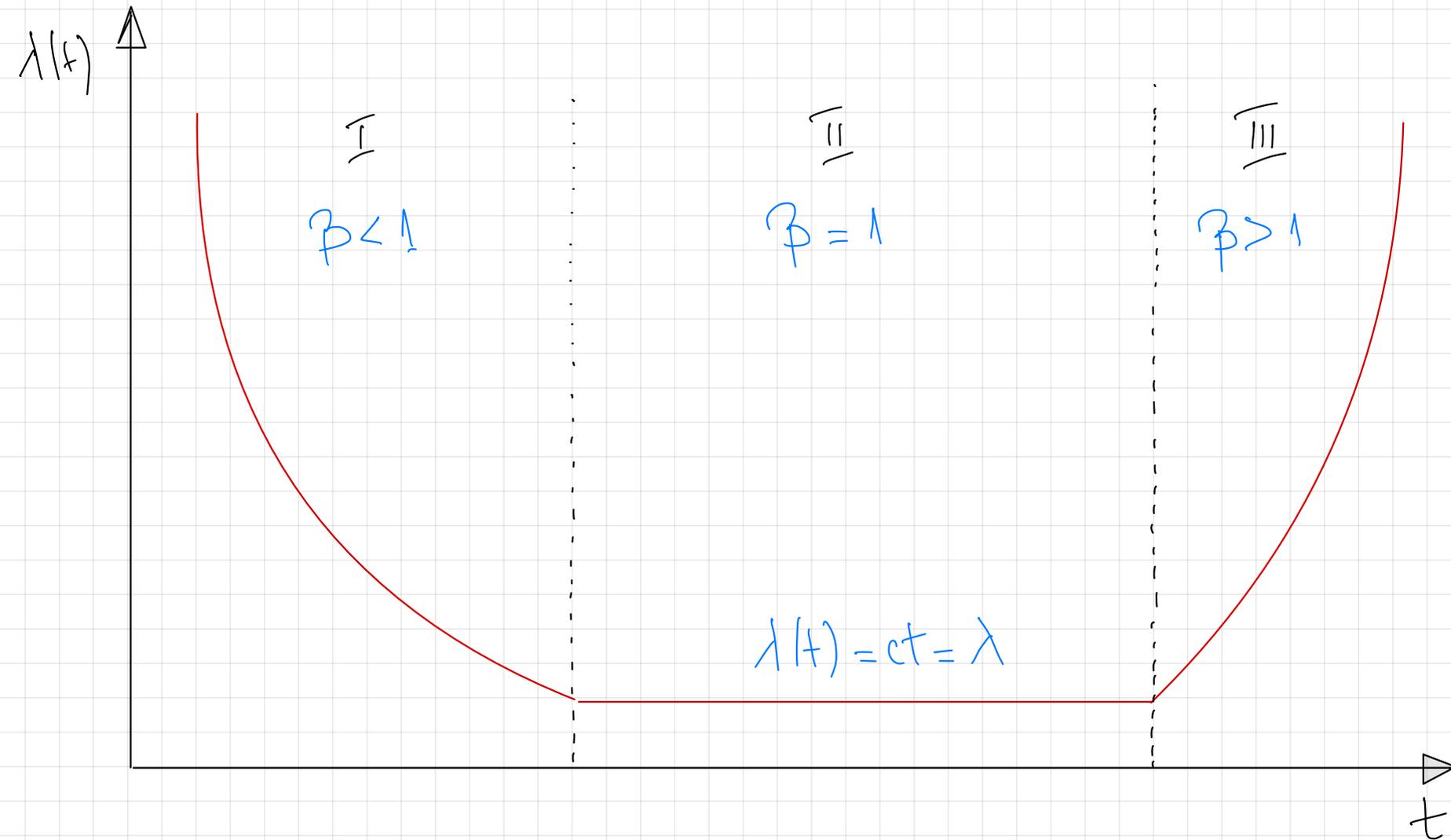
$$\lambda(t) = \frac{- \frac{dR(t)}{dt}}{R(t)} = - \frac{1}{R(t)} \frac{dR(t)}{dt}$$

Aproximare matematică a $\lambda(t)$ - folosim distribuția Weibull

$$f(t) = \lambda \beta t^{\beta-1} e^{-\lambda t^\beta}$$

$$\lambda(t) = \lambda \beta t^{\beta-1}$$

- dacă $\beta < 1 \Rightarrow \lambda(t) \downarrow$ (zona I pe grafic)
- dacă $\beta = 1 \Rightarrow \lambda(t) = ct$ (zona II pe grafic)
- dacă $\beta > 1 \Rightarrow \lambda(t) \uparrow$ (zona III pe grafic)



1. Cazul in care $\lambda(t) = ct = \lambda$ (perioada de viață utilă a produsului)

$$\lambda(t) = \lambda = -\frac{1}{R(t)} \frac{dR(t)}{dt} \Leftrightarrow \lambda dt = -\frac{1}{R(t)} dR(t) \Rightarrow$$

$$\Rightarrow \int \lambda dt = -\int \frac{1}{R(t)} dR(t) \Rightarrow \lambda t + c_1 = -\ln R(t) + c_2 \Rightarrow$$

$$\Rightarrow \ln R(t) = -\lambda t + c \Rightarrow R(t) = e^{-\lambda t + c} = K \cdot e^{-\lambda t}$$

Dacă $t=0 \Rightarrow R(0) = K$ - putem să presupunem că sistemul pornește la momentul zero de timp în stare de funcționare, deci $R(0) = 1$ (fiabilitate 100%) - alegem $K=1 \Rightarrow$

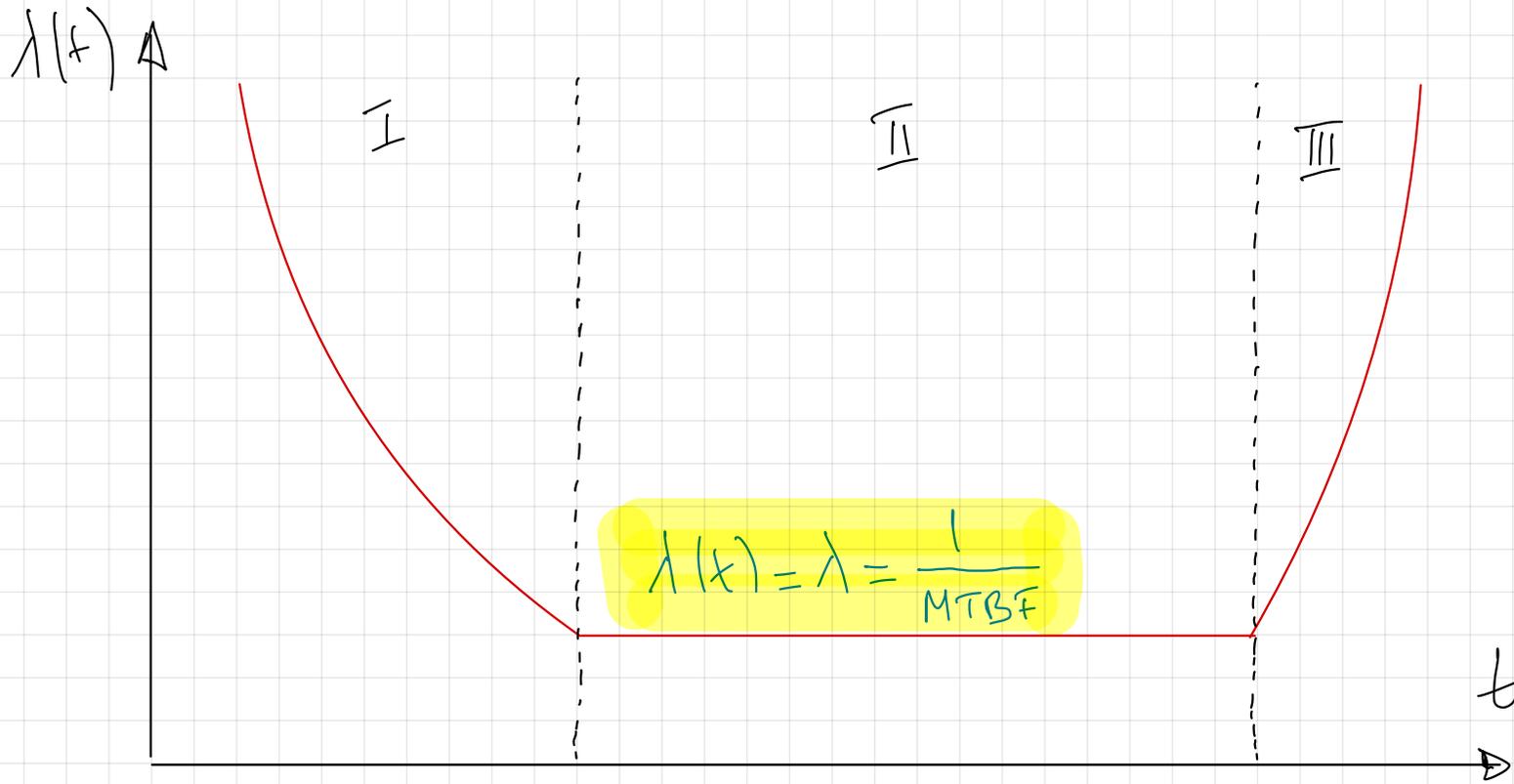
$$R(t) = e^{-\lambda t}$$



$$MTBF = \int_0^{\infty} R(t) dt = \int_0^{\infty} e^{-\lambda t} dt = -\frac{1}{\lambda} e^{-\lambda t} \Big|_0^{\infty} = -\frac{1}{\lambda} (e^{-\lambda \cdot \infty} - e^{-\lambda \cdot 0}) =$$

$$= -\frac{1}{\lambda} (0 - 1) = \frac{1}{\lambda}$$

Deci, pt $\lambda(t) = ct = \lambda$, $MTBF = \frac{1}{\lambda}$



2. Dacă: $\lambda(t) \neq ct$ (pt. zonele montolitare inputul și îmbătrânire)

$$f(t) = \lambda \beta t^{\beta-1} e^{-\lambda t^\beta}$$

$$F(t) = \int_0^t f(z) dz = \int_0^t \lambda \beta z^{\beta-1} e^{-\lambda z^\beta} dz = \dots = 1 - e^{-\lambda t^\beta}$$

dar $R(t) = 1 - F(t) \Rightarrow R(t) = e^{-\lambda t^\beta}$

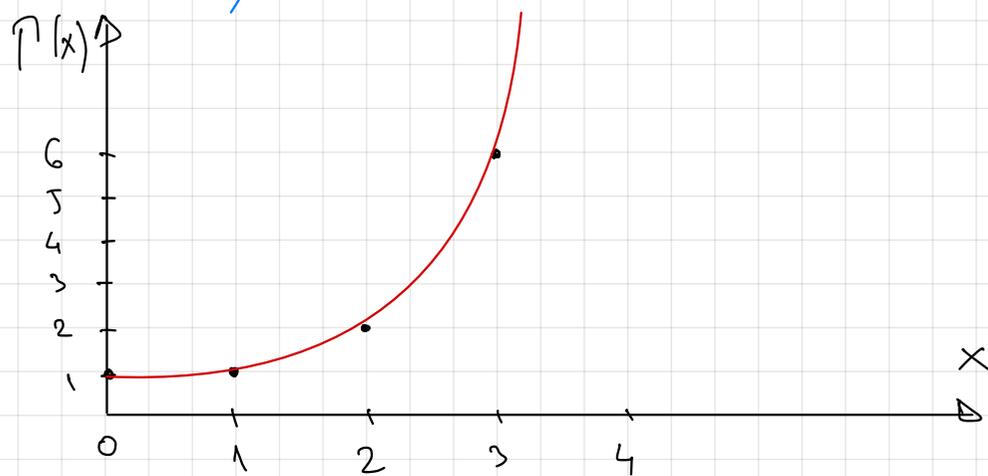
$$MTBF = \int_0^\infty R(t) dt = \int_0^\infty e^{-\lambda t^\beta} dt = \frac{\Gamma(\beta^{-1})}{\beta \lambda^{\beta^{-1}}}, \text{ unde } \Gamma(x) \text{ este funcția}$$

gamma, definită prin: $\Gamma(x) = \int_0^\infty e^{-y} y^{x-1} dy$ și este extensia pt. numere reale a funcției factoriale ($n!$)

$$\Gamma(x) = (-1+x) \Gamma(-1+x)$$

$$\Gamma(x) = \frac{\Gamma(x+1)}{x}$$

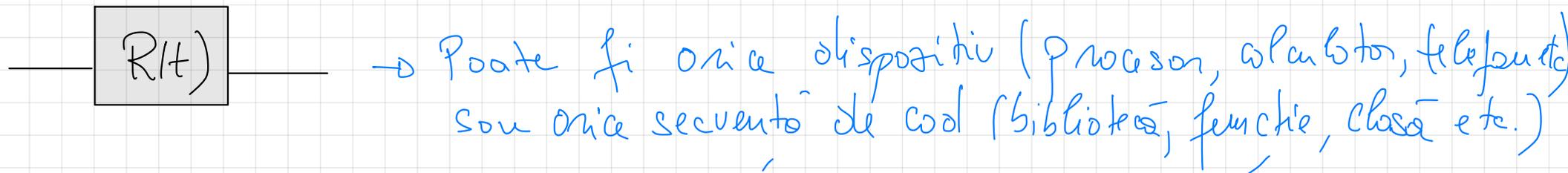
$$\Gamma(0) = \Gamma(1) = 1$$



Estimarea fiabilității

Folosim diagrame pentru a modela și estima fiabilitatea

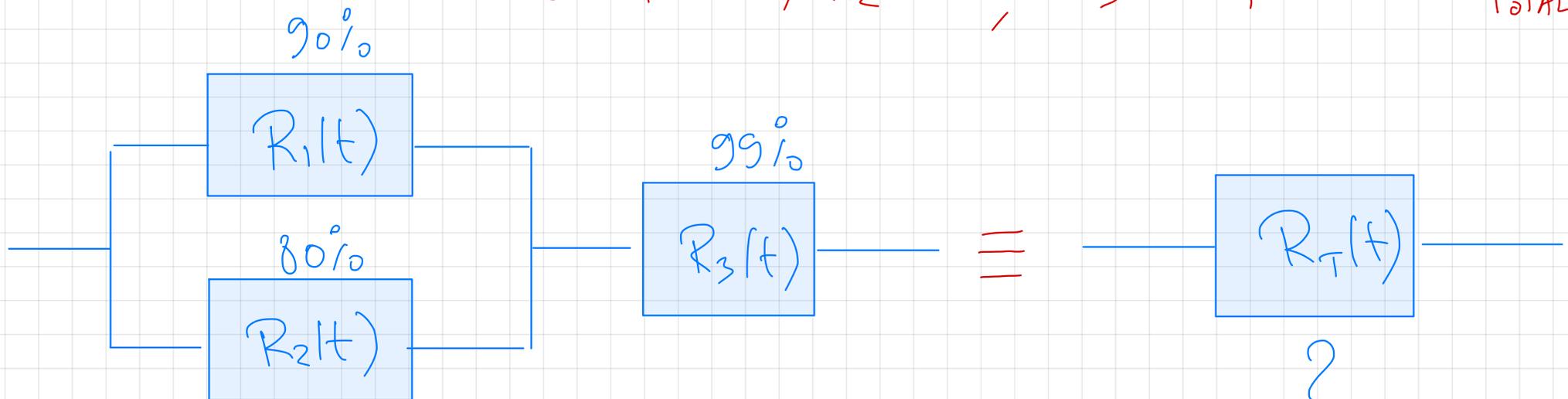
Modularizare: orice sistem, oricât de complex, poate fi modelat fiabilistic ca un modul de fiabilitate dotat $R(t)$:



De ex.: Două procesoare ce accesează aceeași Memorie.

Primul procesor are o fiabilitate $R_1(t)$, al doilea $R_2(t)$ și memoria $R_3(t)$

Dacă $R_1 = 90\%$, $R_2 = 80\%$ și $R_3 = 99\%$, cât este R_{TOTAL} ?



Strukturserie

- n Module in Serie



$$R_s(t) = R_1(t) \cdot R_2(t) \cdot \dots \cdot R_n(t) = \prod_{i=1}^n R_i(t)$$

$$R_s(t) = \prod_{i=1}^n R_i(t)$$

Be ex.:



$$R_s = 0,9 \cdot 0,84 \cdot 0,88 \cdot 0,55 = 0,3659 \approx 0,37 = 37\%$$

Da, ca $R_i(t) = e^{-\lambda_i t}$

$$R_s(t) = \prod_{i=1}^n e^{-\lambda_i t} = e^{-\sum_{i=1}^n \lambda_i t}$$

$$R_s(t) = e^{-\sum_{i=1}^n \lambda_i t}$$

$$\lambda_s = \sum_{i=1}^n \lambda_i$$

Definim $Q_i(t) \rightarrow$ inversul fiabilității

$$Q_i(t) = 1 - R_i(t)$$

$$R_s(t) = \prod_{i=1}^n R_i(t)$$

$$Q_s(t) = 1 - R_s(t) = 1 - \prod_{i=1}^n (1 - Q_i(t))^n$$

R_i este mare $\geq 90\%$ $\Rightarrow Q_i \leq 10\%$ $Q_i, Q_j \approx 0$

$$Q_s(t) = 1 - \left(1 - \sum_{i=1}^n Q_i + \sum_{i=1}^n \sum_{j=1}^n Q_i Q_j - \dots \right) \approx 1 - \left(1 - \sum_{i=1}^n Q_i \right) \Rightarrow$$

$$Q_s(t) = \sum_{i=1}^n Q_i(t)$$

$$MTBF_i = \int_0^{\infty} R_i(t) dt = \int_0^{\infty} e^{-\lambda_i t} dt = \frac{1}{\lambda_i}$$

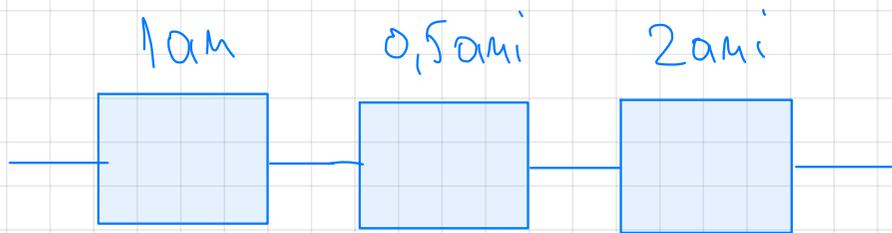
$$MTBF_S = \int_0^{\infty} R_S(t) dt = \int_0^{\infty} e^{-\sum_{i=1}^n \lambda_i t} dt = \int_0^{\infty} e^{-\lambda_S t} dt =$$

$$= -\frac{1}{\lambda_S} e^{-\lambda_S t} \Big|_0^{\infty} = -\frac{1}{\lambda_S} (0 - 1) = \frac{1}{\lambda_S}$$

$$MTBF_S = \frac{1}{\lambda_S}$$

Don $\lambda_i = \frac{1}{MTBF_i} \Rightarrow$

$$MTBF_S = \frac{1}{\lambda_1 + \lambda_2 + \dots + \lambda_n} = \frac{1}{\frac{1}{MTBF_1} + \frac{1}{MTBF_2} + \dots + \frac{1}{MTBF_n}} = \frac{1}{\sum_{i=1}^n \frac{1}{MTBF_i}}$$



$$MTBF_S = \frac{1}{\frac{1}{1} + \frac{1}{0,5} + \frac{1}{2}} = \frac{1}{3 + \frac{1}{2}} = \frac{1}{\frac{7}{2}} = \frac{2}{7} \text{ ami}$$

$$MTBF_S = 0,285 \text{ ami}$$

Dacă avem n module identice în serie, atunci:

$$R_1(t) = R_2(t) = \dots = R_n(t) = R(t) = e^{-\lambda t}$$

$$R_S(t) = \prod_{i=1}^n R_i(t) = \prod_{i=1}^n e^{-\lambda t} = e^{-n\lambda t}$$

$$R_S(t) = e^{-n\lambda t} = R^n(t)$$

$$\lambda_S(t) = \sum_{i=1}^n \lambda_i(t) = \sum_{i=1}^n \lambda = n\lambda$$

$$\lambda_S = n\lambda$$

$$MTBF_S = \frac{1}{\lambda_S} = \frac{1}{n\lambda} = \frac{MTBF}{n}$$

$$MTBF_S = \frac{MTBF}{n}$$

În general, pt structura serie, R total scade, λ total crește și $MTBF$ total scade.

Struktura paralel

$$R_p(t) = ? = 1 - Q_p(t)$$

$$Q_p(t) = Q_1(t) \cdot Q_2(t) \cdot \dots \cdot Q_n(t) = \\ = (1 - R_1(t)) (1 - R_2(t)) \cdot \dots \cdot (1 - R_n(t))$$

$$R_p(t) = 1 - (1 - R_1(t)) (1 - R_2(t)) \cdot \dots \cdot (1 - R_n(t))$$

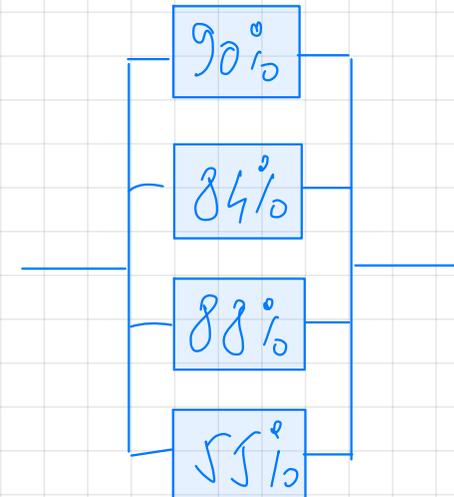
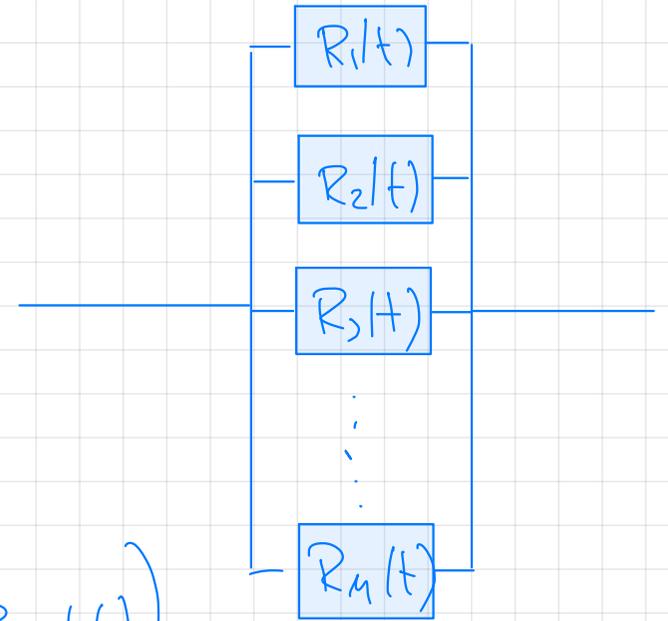
$$R_p(t) = 1 - \prod_{i=1}^n (1 - R_i(t))$$

De ex: patru module in paralel

$$R_p = 1 - (1 - 0,9)(1 - 0,84)(1 - 0,88)(1 - 0,55) =$$

$$= 1 - 0,1 \cdot 0,16 \cdot 0,12 \cdot 0,45 = 0,9991$$

$$R_p = 99,91\%$$



$$R_i(t) = e^{-\lambda_i t} \Rightarrow R_p(t) = 1 - \prod_{i=1}^n (1 - e^{-\lambda_i t})$$

$$R_p(t) = 1 - (1 - e^{-\lambda_1 t})(1 - e^{-\lambda_2 t}) \dots (1 - e^{-\lambda_n t}) \stackrel{?}{=} e^{-\lambda_p t}$$

Presupunem ca $R_1(t) = R_2(t) = \dots = R_n(t) = e^{-\lambda t}$

$$R_p(t) = 1 - (1 - e^{-\lambda t})^n$$

$$\lambda_p(t) = \frac{f_p(t)}{R_p(t)} = \frac{-\frac{dR_p(t)}{dt}}{R_p(t)}$$

$$\frac{dR_p(t)}{dt} = -n(1 - e^{-\lambda t})^{n-1} (1 - e^{-\lambda t})' = -n\lambda e^{-\lambda t} (1 - e^{-\lambda t})^{n-1}$$

$$\lambda_p(t) = n\lambda \frac{e^{-\lambda t} (1 - e^{-\lambda t})^{n-1}}{1 - (1 - e^{-\lambda t})^n}$$

- nu este constantă!

$$\lambda_{P, \text{steady state}} = \lim_{t \rightarrow \infty} \lambda_P(t) = \lim_{t \rightarrow \infty} m \lambda \frac{e^{-\lambda t} (1 - e^{-\lambda t})^{m-1}}{1 - (1 - e^{-\lambda t})^m} =$$

$$e^{-\lambda t} = x, \text{ odc\u0107a } t \rightarrow \infty, \text{ atunci } x \rightarrow 0$$

$$= \lim_{x \rightarrow 0} m \lambda \frac{x(1-x)^{m-1}}{1 - (1-x)^m} \stackrel{\text{l'Hospital}}{=} m \lambda \lim_{x \rightarrow 0} \frac{(x(1-x)^{m-1})'}{(1 - (1-x)^m)'} =$$

$$= m \lambda \lim_{x \rightarrow 0} \frac{(1-x)^{m-1} - x(m-1)(1-x)^{m-2}}{+ m(1-x)^{m-1}} = m \lambda \lim_{x \rightarrow 0} \frac{1-x - x(m-1)}{m(1-x)}$$

$$= m \lambda \cdot \frac{1}{m} = \lambda$$

$$\lambda_P = \lambda, \text{ la steady-state}$$

$$MTBF_P = \int_0^{\infty} R_p(t) dt = \int_0^{\infty} \left[1 - \prod_{i=1}^n (1 - R_i(t)) \right] dt$$

Doi modulele sunt identice $R(t) = e^{-\lambda t}$

$$MTBF_P = \int_0^{\infty} \left(1 - (1 - e^{-\lambda t})^n \right) dt$$

$$(1 - e^{-\lambda t})^n = 1 - n e^{-\lambda t} + C_n^2 e^{-2\lambda t} - C_n^3 e^{-3\lambda t} + \dots + (-1)^n e^{-n\lambda t}$$

$$MTBF_P = \int_0^{\infty} \left(n e^{-\lambda t} - \frac{n(n-1)}{2} e^{-2\lambda t} + \dots + (-1)^{n+1} e^{-n\lambda t} \right) dt =$$

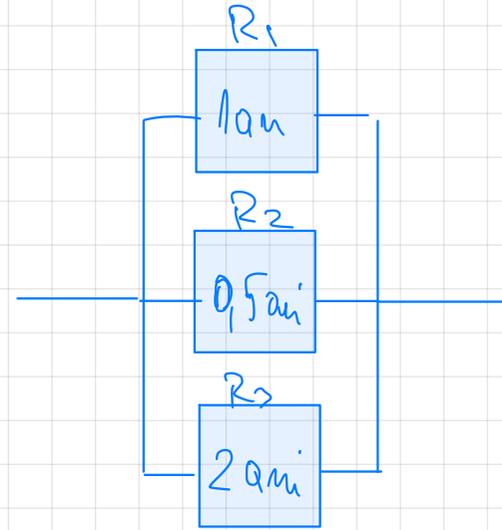
$$= n \frac{1}{\lambda} - \frac{n(n-1)}{2} \frac{1}{2\lambda} + \dots + (-1)^{n+1} \frac{1}{n\lambda} = \frac{1}{\lambda} \left(n - \frac{n(n-1)}{2} + \dots + (-1)^{n+1} \frac{1}{n} \right)$$

$$= \frac{1}{\lambda} \sum_{i=1}^n C_n^i \frac{1}{i} (-1)^{i+1}$$

In general, R total creste, λ total rămâne constant și $MTBF$ total creste

De ex.: trei module in paralel, $MTBF_1 = 1 \text{ an}$, $MTBF_2 = 0,5 \text{ ani}$, $MTBF_3 = 2 \text{ ani}$

$$MTBF_P = \int_0^{\infty} R_P(t) dt = \int_0^{\infty} (1 - (1-R_1)(1-R_2)(1-R_3)) dt$$



$$(1-R_1)(1-R_2)(1-R_3) = (1-R_2-R_1+R_1R_2)(1-R_3) =$$

$$= 1 - R_1 - R_2 + R_1R_2 - R_3 + R_2R_3 + R_1R_3 - R_1R_2R_3$$

$$\int_0^{\infty} (R_1 + R_2 + R_3 - R_1R_2 - R_2R_3 - R_1R_3 + R_1R_2R_3) dt$$

$$\int_0^{\infty} (e^{-\lambda_1 t} + e^{-\lambda_2 t} + e^{-\lambda_3 t} - e^{-(\lambda_1+\lambda_2)t} - e^{-(\lambda_2+\lambda_3)t} - e^{-(\lambda_1+\lambda_3)t} + e^{-(\lambda_1+\lambda_2+\lambda_3)t}) dt$$

$$= \frac{1}{\lambda_1} + \frac{1}{\lambda_2} + \frac{1}{\lambda_3} - \frac{1}{\lambda_1+\lambda_2} - \frac{1}{\lambda_2+\lambda_3} - \frac{1}{\lambda_1+\lambda_3} + \frac{1}{\lambda_1+\lambda_2+\lambda_3} =$$

$$= MTBF_1 + MTBF_2 + MTBF_3 - \frac{1}{\frac{1}{MTBF_1} + \frac{1}{MTBF_2}} - \frac{1}{1+2} - \frac{1}{2+\frac{1}{2}} - \frac{1}{1+\frac{1}{2}}$$

$$+ \frac{1}{1+2+\frac{1}{2}} = 3,5 - \frac{1}{3} - \frac{2}{5} - \frac{2}{3} + \frac{2}{7} = 2,385 \text{ ani}$$

Conclusi:

n module identica in Serie

n module identica in Parallelo

R_{TOTAL}

$$R_S(t) = R^n(t) = e^{-n\lambda t}$$

$$R_P(t) = 1 - (1 - R(t))^n = 1 - (1 - e^{-\lambda t})^n$$

λ_{TOTAL}

$$\lambda_S = n\lambda$$

$$\lambda_P(t) = n\lambda \frac{e^{-\lambda t} (1 - e^{-\lambda t})^{n-1}}{1 - (1 - e^{-\lambda t})^n}, \quad \lambda_{P, \text{stabile}} = \lambda$$

$MTBF_{TOTAL}$

$$MTBF_S = \frac{MTBF}{n}$$

$$MTBF_P = MTBF \cdot \sum_{i=1}^n (-1)^{i+1} \cdot C_n^i \frac{1}{i}$$