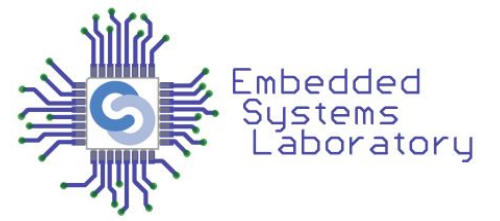


Internet of Things

Lecture 4 - 6LoWPAN & RPL



6LoWPAN

IPv4

- Exhausted in 2011
- 32-bit address

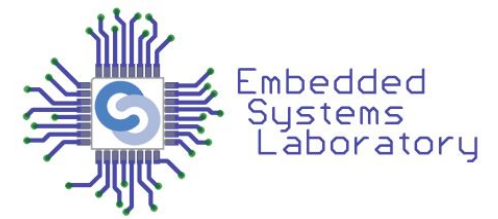
IPv6

- 128-bit addresses
- Fit for large IoT networks
- Not enough resources on low power devices

6LoWPAN

- Adaptation layer
- Header compression
- Fragmentation

6LoWPAN Standard



- Internet Engineering Task Force (IETF)
 - TCP, UDP, HTTP, CoAP, etc.
- RFC 4944 - first 6LoWPAN standard
- RFC 6282 - header compression
- RFC 6775 - neighbor discovery
- Over IEEE 802.15.4
- Adapted to work with other low-power technologies
 - Bluetooth Smart
 - Wi-Fi low-power

Network Architecture

- 6LoWPAN mesh network
- Edge router (6LBR)
 - Generates & manages 6LoWPAN network
 - Transfers between 6LoWPAN nodes and the Internet
 - Transfers between 6LoWPAN nodes
 - Network level
- Routers (6LR)
 - Routes traffic from one node to another
- Hosts (6LN)
 - Do not route traffic
 - Low power nodes

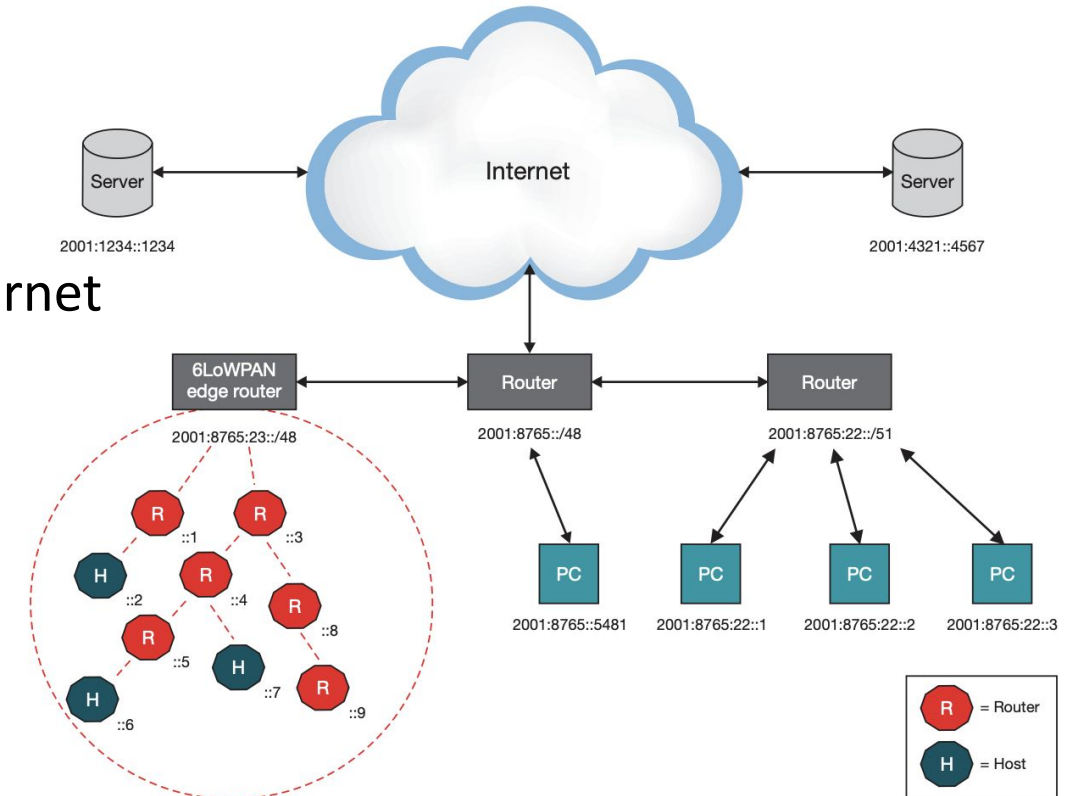


Figure 1. An example of an IPv6 network with a 6LoWPAN mesh network

Source: Olsson, Jonas. "6LoWPAN demystified." Texas Instruments 13 (2014).

Networking Stack

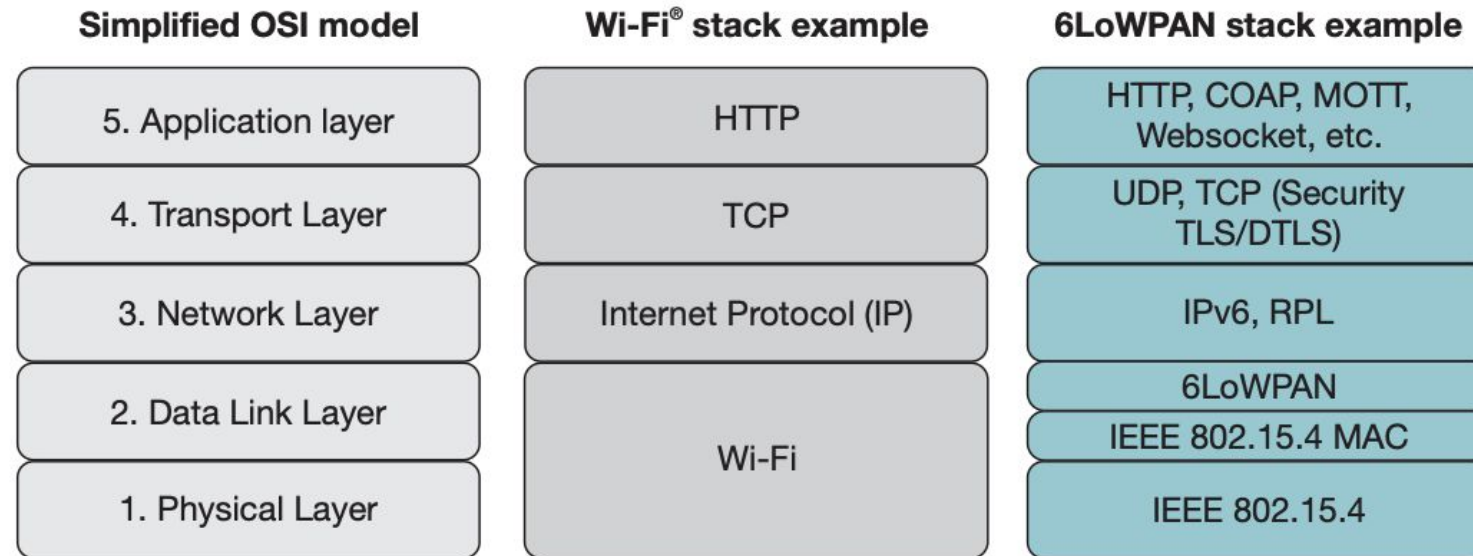
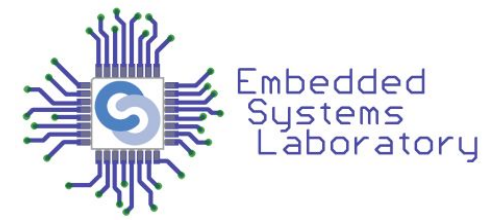


Figure 2. The OSI model, a Wi-Fi[®] stack example and the 6LoWPAN stack

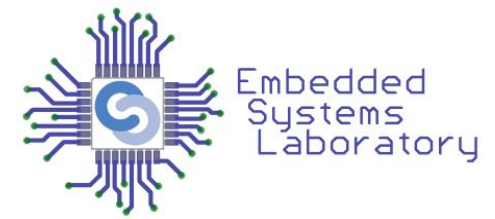
Image source: Olsson, Jonas. "6LoWPAN demystified." Texas Instruments 13 (2014).

IPv6 over IEEE 802.15.4



- IPv6
 - MTU is 1280 bytes
 - Reflects technology advancement
- 802.15.4
 - Low power, low cost devices
 - Frame size is 127 bytes
 - Maximum bandwidth 250 Kbps
 - MAC addresses on 64 bits or 16 bits
 - Minimize header overhead & memory consumption

IPv6 over IEEE 802.15.4



- Main challenges for using IPv6 over 802.15.4
- IPv6 has minimum MTU 10 times larger
 - IPv6 - 40 bytes headers, TCP - 20 bytes, UDP - 8 bytes
 - Solution: Fragmentation & header compression
- Low power and lossy networks
 - Interferences, unstable links, packet loss
 - Solution: Adaptive and responsive network layer

- 6LoWPAN Working Group from IETF => RFC 6282
 - Encapsulation of IPv6 packet into 802.15.4 frame
- Header compression
 - Stateless or context-based compression
 - The elimination of header fields that can be derived from other headers
 - Same network prefix
 - Determine IPv6 addresses and field sizes

- Fragmentation
 - Fragment data to fit in 802.15.4 frames
- Stateless auto-configuration
 - 6LoWPAN nodes generate their own addresses
 - Duplicate address detection (DAD)

Header Compression

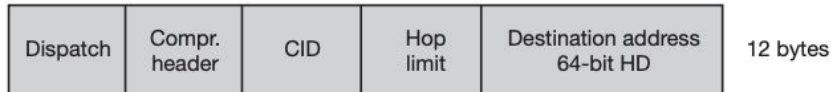
IPv6 header



1. Compressed header, FE80::CAFE:00FF:FE00:0100 → FE80::CAFE:00FF:FE00:0200



2. Compressed header, 2001::DEC4:E3A1:FE24:9600 → 2001::4455:84C6:39BB:A2DD



3. Compressed header, 2001::DEC4:E3A1:FE24:9600 → 2001::4455:84C6:39BB:A2DD



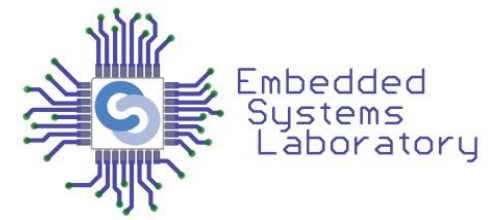
- Stateless & shared-context compression
- IPv6 addr = network prefix + interface ID
- Interface ID - derived from MAC address

1. between nodes from the local network
 - exclude both link-local addresses
 - useful for routing protocols
2. destination is external but prefix is known
 - exclude src addr & dest prefix
3. destination is external and prefix is not known
 - exclude interface ID

Figure 3. 6LoWPAN IPv6 header compression examples

Image source: Olsson, Jonas. "6LoWPAN demystified." Texas Instruments 13 (2014).

Fragmentation



- Packets are divided into smaller segments
- Additional information in the headers for the reassembly
- **mesh-under** routing:
 - packets are reassembled at the destination
 - quick routing of fragments, reduced delay
 - if a single fragment is lost the whole packet must be retransmitted
- **route-over** routing:
 - packet is reconstructed at each hop
 - hops are devices with more resources
- Avoid fragmentation - reduced payload + header compression

Stacked headers

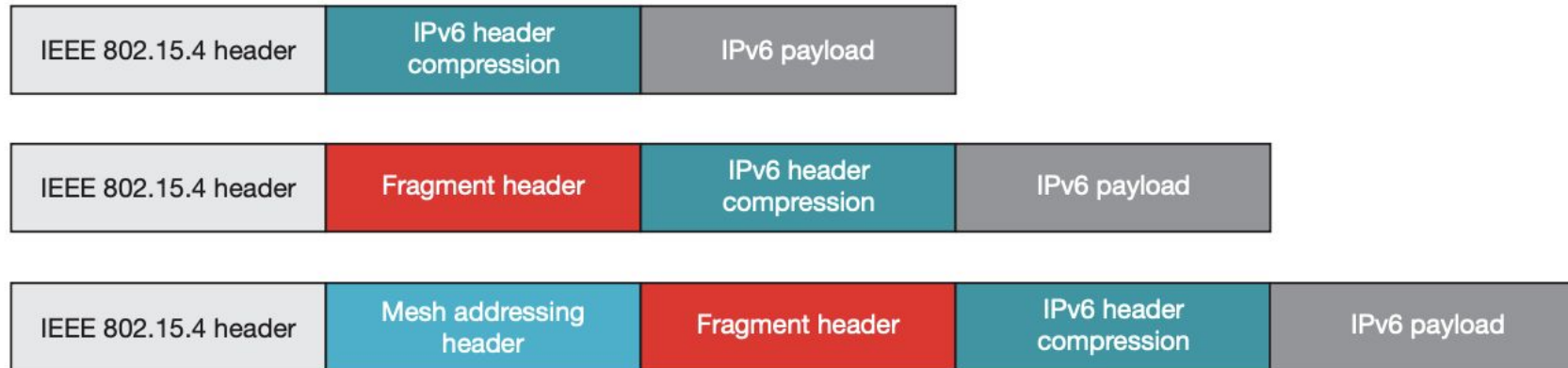


Figure 4. 6LoWPAN stacked headers

Image source: Olsson, Jonas. "6LoWPAN demystified." Texas Instruments 13 (2014).

3 types of sub-headers:

- mesh addressing
 - multi-hop topology
 - forward packets
- fragmentation
 - identify fragments
- header compression

- Fragment header
 - IPv6 packet is bigger than 802.15.4 frame
 - Fragmentation is needed
 - 3 fields:
 - datagram size
 - whole payload size
 - datagram tag
 - identify a set of fragments of the same payload
 - datagram offset
 - location of fragment in payload

- Mesh addressing header
 - Used in multi-hop topologies for routing
 - 3 fields:
 - hop limit
 - decremented at each hop
 - drop packet when hop limit reaches 0
 - source address
 - destination address
 - 802.15.4 MAC addresses
 - 16 or 64 bits

- Mesh-under
 - Uses 802.15.4 MAC addresses to forward packets
 - Forwarding is done at link layer
 - IP subnet
 - A single broadcast domain
 - Router IP = edge router
 - All messages are sent to all nodes
 - Fit for small networks

- Route-over
 - Works at network layer
 - Each node is an IP router
 - Each node implements all functionality (including DAD)
 - Recommended for large networks
 - RPL is a route-over protocol

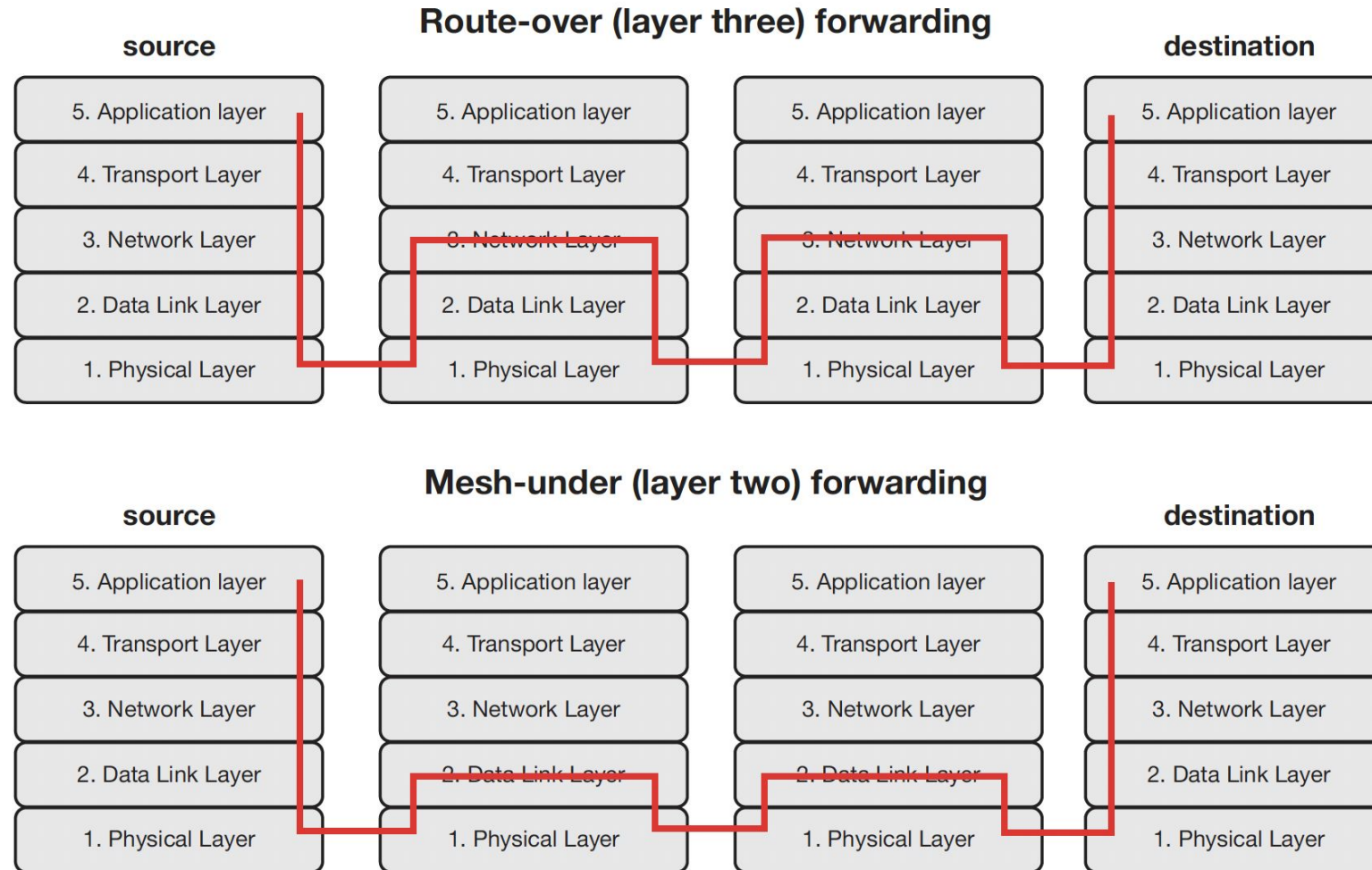
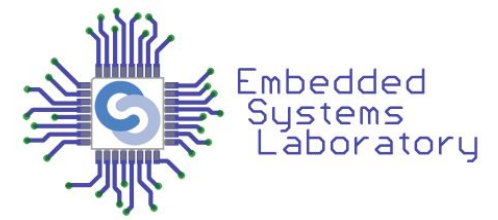


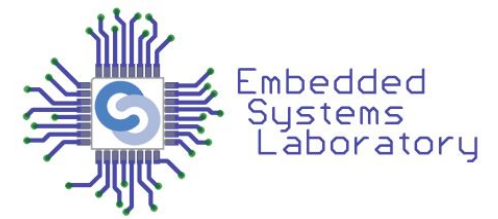
Figure 5. Mesh-under and route-over packet forwarding

Neighbor Discovery



- Neighbor Discovery Protocol (NDP)
 - used for discovering neighbor devices
 - maintain information about available devices
 - configure default routes
 - propagate configuration parameters
- 4 types of messages:
 - Router solicitation (RS)
 - Router advertisement (RA)
 - Neighbor solicitation (NS)
 - Neighbor advertisement (NA)

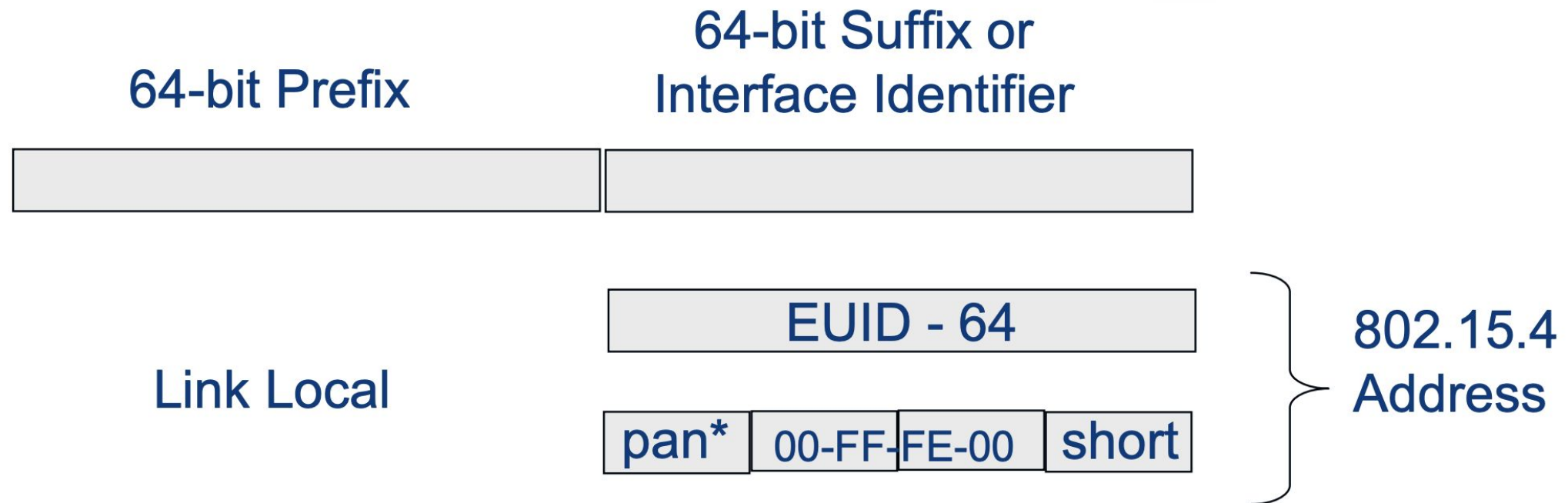
Neighbor Discovery



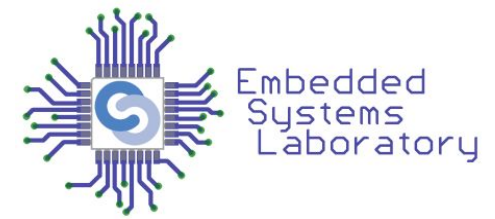
- RS/RA
 - discover routers
 - find network prefix
- NS/NA
 - find duplicate addresses (DAD)
 - node generates link-local address and sends NS for verification
 - if it receives NA with duplicate flag, the address is not unique
 - finding neighbours
- Using these 4 messages a node can generate a unique address

Link-local address

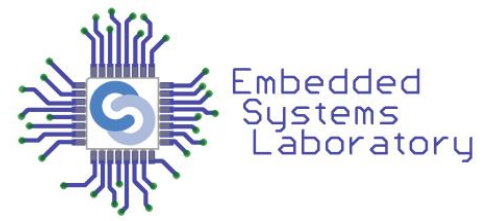
- Auto-generate IPv6 address, no need for DHCP
- Link-local address derived from 802.15.4 address (64 or 16 bits)
- Link-local prefix is FE80::/64



Advantages of auto-configuration



- Mesh-under routing
 - local-link address is sufficient to communicate within the local 6LoWPAN network
 - cannot be used for communicating outside the local network
- Router-over routing
 - local-link address for communicating with direct neighbors
 - for multi-hop communication it needs routable address
- Advantage of deriving the IPv6 address from 802.15.4 MAC address
 - eliminate some fields to compress headers
- Same prefix in the network
 - also good for header compression
 - prefix is discovered through RS/RA messages



RPL

- RPL = Routing Protocol for Low power and lossy networks
- Defined by IETF in RFC 6550
- IETF ROLL working group
- IP smart object networks / Low-power and lossy networks
- “route-over” protocol for 6LoWPAN networks
- Distance-vector & source routing protocol
- Communication:
 - multipoint-to-point
 - point-to-multipoint
 - point-to-point

- Directed Acyclic Graph (DAG)
 - similar to a tree
 - a node can associate to multiple parents
- Destination-Oriented DAGs (DODAGs),
 - Sink node/gateway - root of the DAG
- RPL instance = one or more DODAGs
 - *RPLInstanceID* identifies the instance
 - An RPL network may have multiple instances
 - A node may belong to several instances, but only to one DODAG in each instance

RPL Topology Example

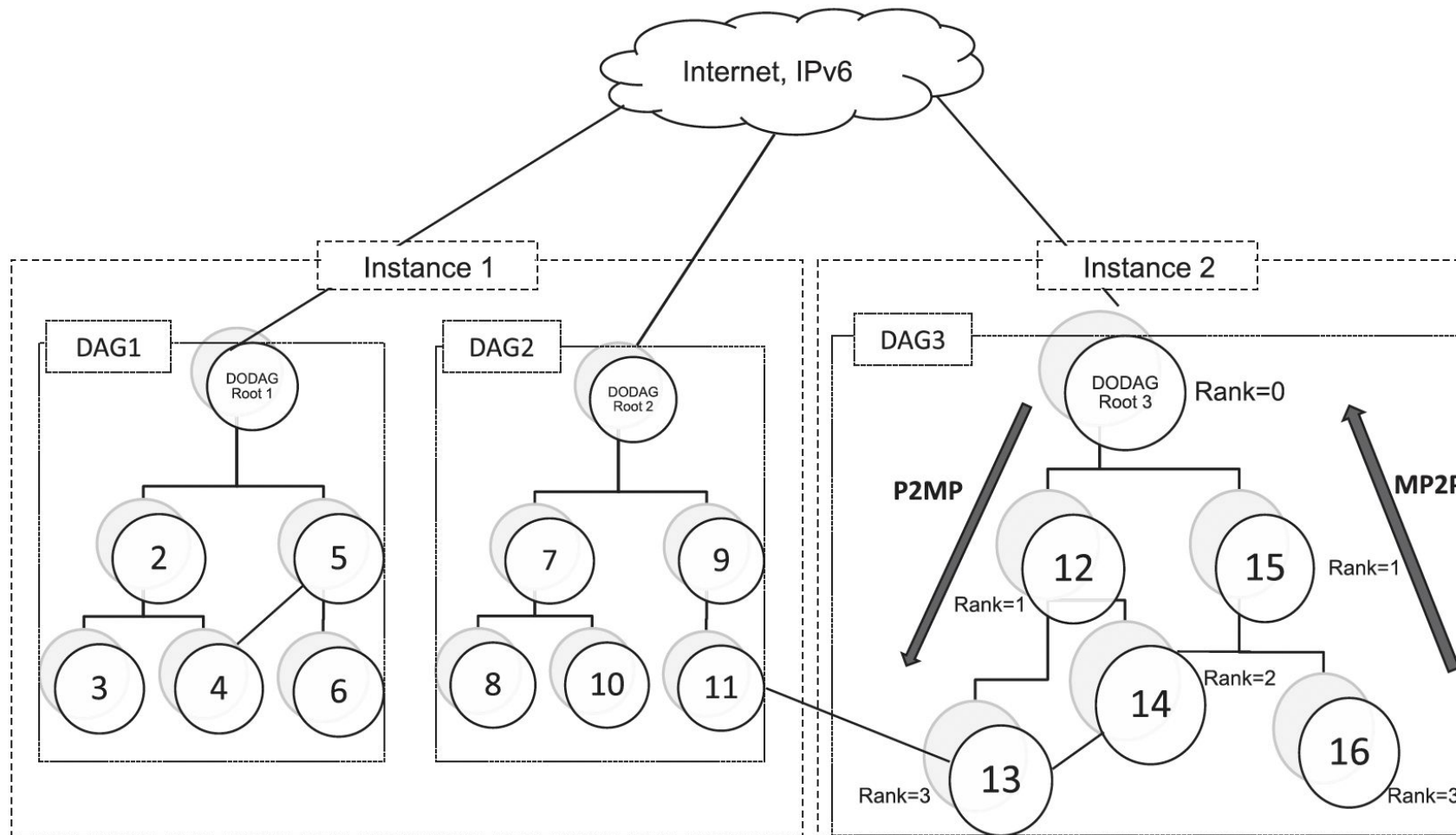


Image source: Gaddour, Olfa, and Anis Koubâa. "RPL in a nutshell: A survey." *Computer Networks* 56.14 (2012): 3163-3178.

Fig. 1. A RPL network with three DODAGs in two instances.

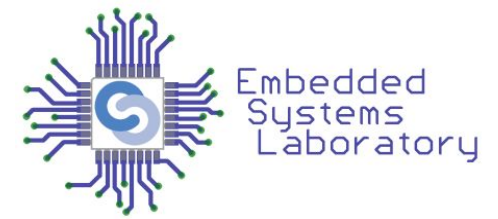
- Combines both hierarchical and mesh topologies
- By default it's a hierarchical topology
 - child-parent relationship
 - sending packets to the preferred parent
- Allows routing through siblings or children
 - when the parent is not available
- Flexibility

- Auto-configuration
 - Compatible with 6LoWPAN
 - Neighbor Discovery
 - Build paths to certain destinations
- Self-healing
 - Nodes may become unavailable
 - Choose more than one parent for redundancy
 - Network adapts to changes

- Loop avoidance & detection
 - DAG is acyclic
 - a node has a higher rank than its parents
 - detects loops and fixes the network
- Independence from link layer protocols
 - runs over multiple link-layer technologies
- Multiple edge routers
 - Multiple DODAGs
 - Each DAG root is an edge router

- 3 types of nodes:
 - LBR (Low Power and Lossy Border Router)
 - DODAG root
 - create DAG
 - gateway / edge router
 - Router
 - route packets
 - generate data packets
 - attach to existing DAG
 - Host
 - generate data packets

Network Model



- Every node has a **rank**
 - integer value
 - position relative to other nodes
 - increases in downstream direction
 - used to detect and avoid loops
- Node can be associated with parent or sibling
 - Rank is used to differentiate between parent & sibling
 - List of potential parents and siblings in case the current parent fails
 - A single parent will be the preferred one (based on a metric)

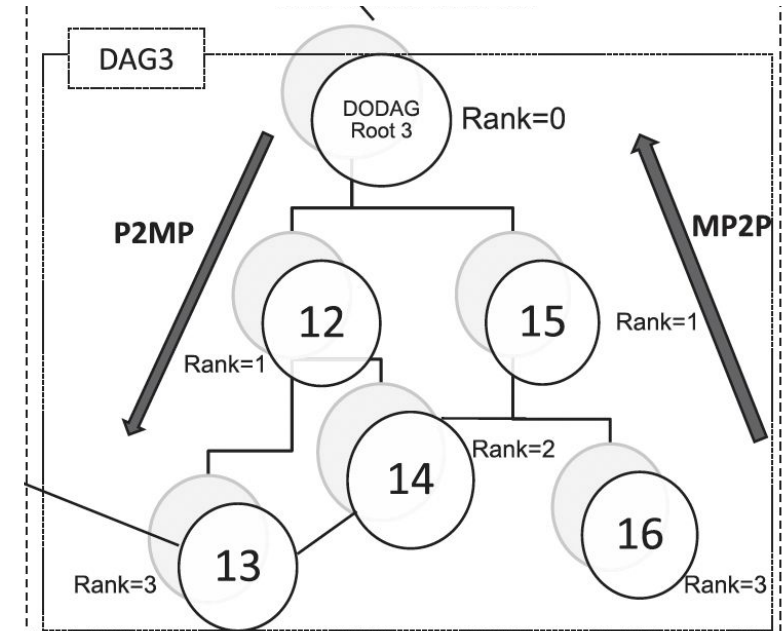


Image source: Gaddour, Olfa, and Anis Koubâa. "RPL in a nutshell: A survey." *Computer Networks* 56.14 (2012): 3163-3178.

DODAG Construction

- Creating a DODAG is based on 2 operations:
 - broadcast **DIO** control messages
 - sent initially by the root
 - each node computes rank
 - adds parents to list
 - forwards DIO
 - upward routes
 - unicast **DAO** control messages
 - generated by nodes
 - set to preferred parent
 - downward routes

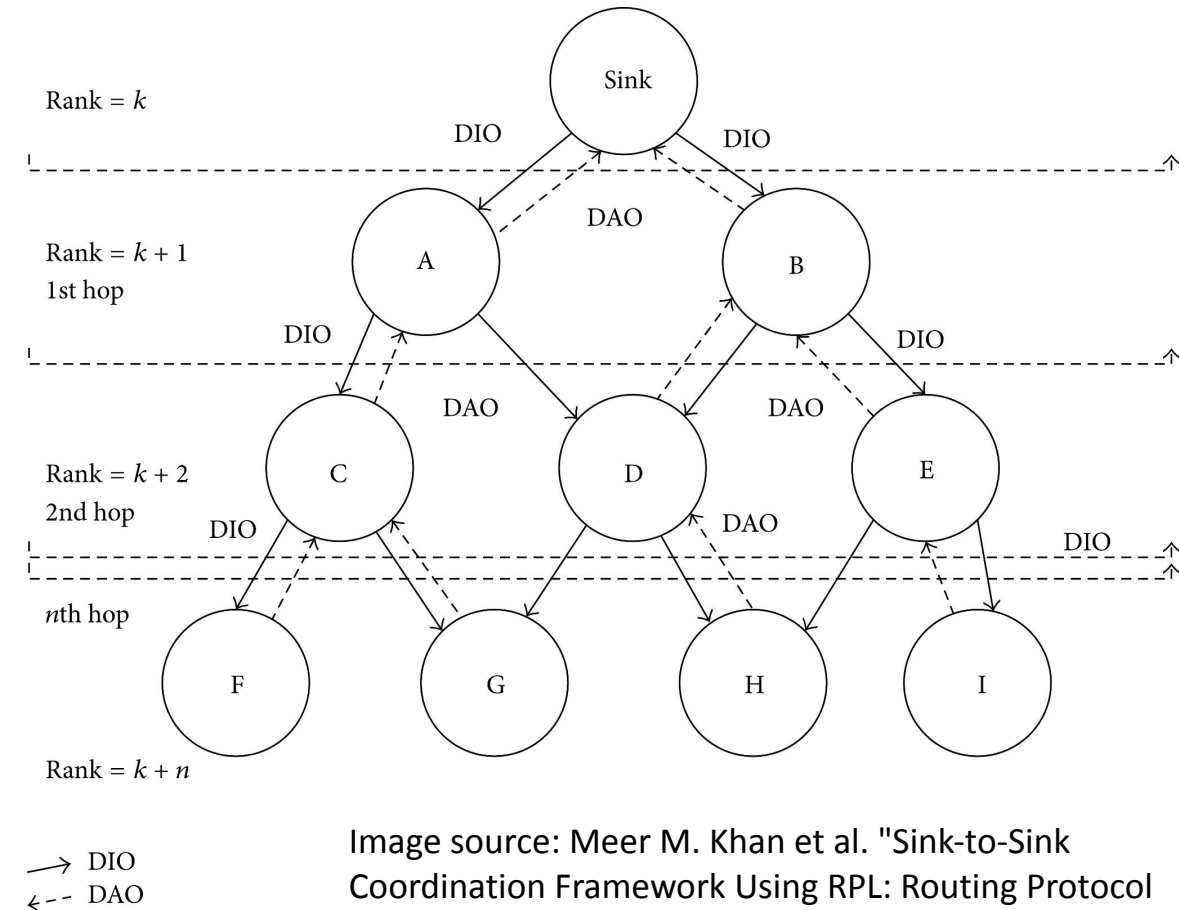


Image source: Meer M. Khan et al. "Sink-to-Sink Coordination Framework Using RPL: Routing Protocol for Low Power and Lossy Networks", Journal of Sensors, vol. 2016, Article ID 2635429, 11 pages, 2016

DODAG Construction

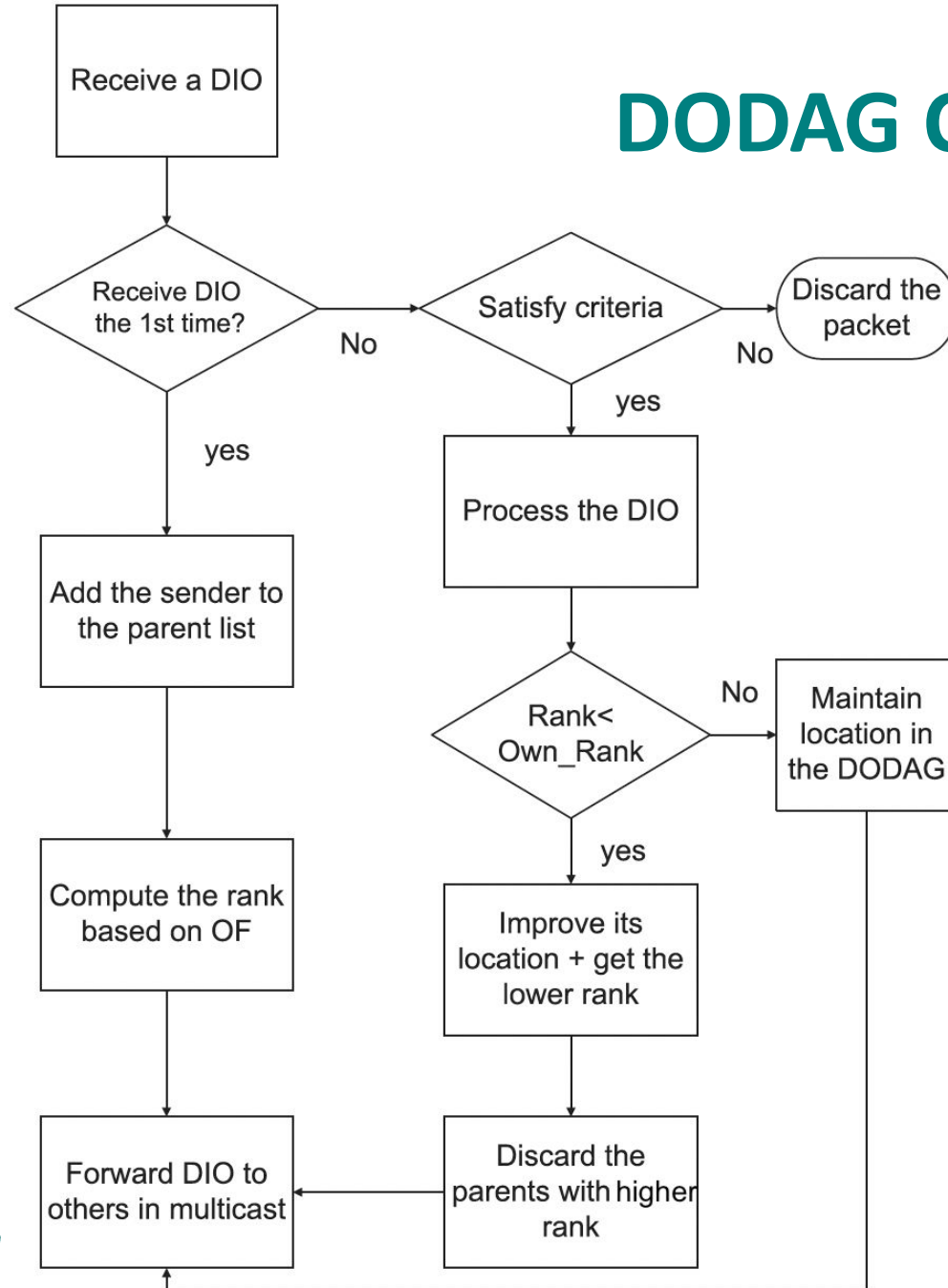
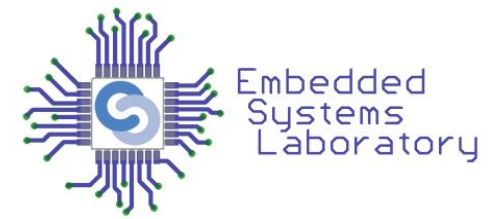


Fig. 5. The operation of a router in a DODAG.

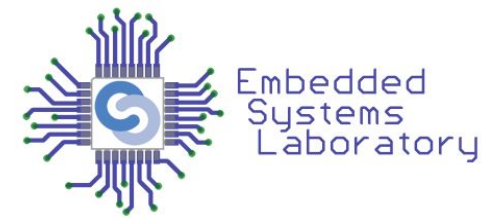
- DIO broadcast (DODAG ID, rank, objective function)
- Node receives first DIO
 - Adds sender to parent list
 - Computes its rank based on OF
 - Forwards the message with updated rank
 - Chooses default parent (default route)
- Node receives another DIO
 - Computes rank
 - If rank is lower, update rank
 - Discard parents with higher rank

Modes of Operation



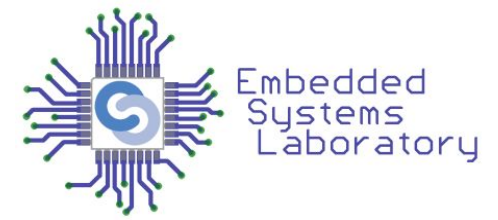
- DIO messages contain a mode of operation flag
- Mode of operation is different than zero
 - downward routes must be obtained
 - each node sends a DAO to parent
 - each parent adds its address and forwards DAO to its parent
 - until it reaches the root
 - full route between the root and each node

Modes of Operation



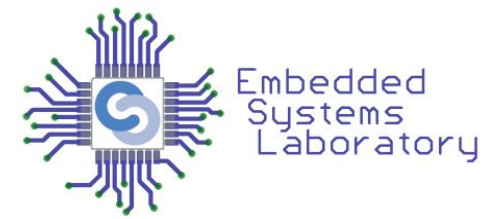
- Two modes to build downward routes: storing & non-storing
- Storing mode:
 - Parents aggregate DAO messages from all children then forwards to parent
 - Maintains a routing table & a neighbour table
 - Routing table includes routes to certain destinations
 - Neighbour table includes direct neighbours
 - Needs more storing space

Modes of Operation



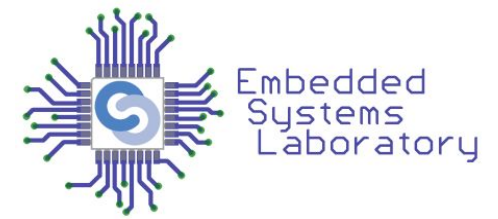
- Non-storing mode:
 - Parents add their addresses and forward DAO without storing the message
 - Only the root maintains a routing table
 - Source-routing
 - Packet includes the complete route (hops) to destination
 - A packet sent from one node to another goes through the root

Bibliography



- RFC 4944: <https://datatracker.ietf.org/doc/html/rfc4944>
- RFC 6282: <https://datatracker.ietf.org/doc/html/rfc6282>
- RFC 6775: <https://datatracker.ietf.org/doc/html/rfc6775>
- Olsson, Jonas. "6LoWPAN demystified." Texas Instruments 13 (2014).
- RFC 6550: <https://datatracker.ietf.org/doc/html/rfc6550>
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Keywords



- IPv6
- 6LoWPAN
- Header compression
- Fragmentation
- Routing
- RPL
- DODAG
- Neighbor discovery
- Auto-configuration